

# Chains of theories

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This is work with Özcan Kasal. There is some parallel work by Alice Medvedev (presented in Oléron, June, 2011) concerning ACFA.

Suppose

$$T_0 \subseteq T_1 \subseteq T_2 \subseteq \dots,$$

all theories, closed under entailment, so their signatures also form a chain:

$$\mathcal{S}_0 \subseteq \mathcal{S}_1 \subseteq \mathcal{S}_2 \subseteq \dots$$

In one example of interest,  $T_m$  is  $m$ -DF, the theory of fields with  $m$  commuting derivations  $\partial_0, \dots, \partial_{m-1}$ ; their union is  $\omega$ -DF.

In general, what properties are preserved in  $\bigcup_{k \in \omega} T_k$ ? Compare:

**Theorem** (Chang, Łoś–Suszko). *For a fixed theory  $T$ , the following are equivalent:*

1.  $T$  is  $\forall\exists$ -axiomatizable.
2.  $\text{Mod}(T)$  is closed under taking unions of chains

$$\mathfrak{A}_0 \subseteq \mathfrak{A}_1 \subseteq \mathfrak{A}_2 \subseteq \dots$$

Again, if  $T_0 \subseteq T_1 \subseteq T_2 \subseteq \dots$ , then among possible properties of the theories  $T_k$ ,

- 1) Preserved by  $\bigcup_{k \in \omega} T_k$  are: (a) consistency, (b) completeness, (c) quantifier elimination, (d) model-completeness, (e) stability, (f) ...;
- 2) not preserved (but this is not obvious) are: companionability,  $\omega$ -stability, superstability, ...

(a) Consistency is preserved, by compactness.

(b) Completeness is preserved, because every sentence of the union  $\bigcup_{k \in \omega} \mathcal{S}_k$  is a sentence of some  $\mathcal{S}_k$ .

(c) Likewise for quantifier-elimination.

(d) **Model-completeness** of a theory  $T$  may be usually remembered as

$$\mathfrak{A} \subseteq \mathfrak{B} \implies \mathfrak{A} \preceq \mathfrak{B}$$

within  $\text{Mod}(T)$ . Equivalently (the theory axiomatized by)

$$T \cup \text{diag}(\mathfrak{A})$$

is always complete when  $\mathfrak{A} \models T$ , where

$$\text{diag}(\mathfrak{A}) = \{\sigma \in \text{Th}(\mathfrak{A}_A) : \sigma \text{ is quantifier-free}\}$$

the theory of the structures in which  $\mathfrak{A}$  embeds. A sufficient (and obviously necessary) condition is (Abraham) **Robinson's Condition**,

$$\mathfrak{A} \subseteq \mathfrak{B} \implies \mathfrak{A} \preceq_1 \mathfrak{B}$$

where the conclusion means every quantifier-free formula over  $\mathfrak{A}$  soluble in  $\mathfrak{B}$  is soluble in  $\mathfrak{A}$ . Robinson's Condition is, equivalently,

$$\mathfrak{A} \models_{\text{ec}} T$$

—every model of  $T$  is an **existentially closed** model. For this it is sufficient (and in fact necessary) that  $T$  admit quantifier-elimination down to  $\exists$  formulas. Therefore model-completeness is preserved in unions of chains.

(e) A *complete* theory  $T$  is

- **$\kappa$ -stable**, if  $\kappa \geq |T|$  and

$$|A| \leq \kappa \implies |S(A)| \leq \kappa$$

for all parameter-sets  $A$  of models of  $T$ ;

- **superstable**, if  $\kappa$ -stable for  $\kappa$  large enough;
- **stable**, if  $\kappa$ -stable for some  $\kappa$ .

When  $|T| = \omega$ , then

- superstability implies  $\kappa$ -stability when  $\kappa \geq 2^\omega$ ;
- stability implies  $\kappa$ -stability when  $\kappa = \kappa^\omega$ .

(Note that if  $\text{cof}(\kappa) = \omega$ , as when  $\kappa = \aleph_\omega$ , then  $\kappa < \kappa^\omega$ .)

In fact *instability* of  $T$  is equivalent to the presence of a formula  $\varphi(\vec{x}, \vec{y})$  defining an infinite linear order in some model of  $T$ , so that, for all  $n$  in  $\omega$ ,

$$T \vdash \exists(\vec{x}_0, \dots, \vec{x}_n) \left( \bigwedge_{0 \leq i \leq j \leq n} \varphi(\vec{x}_i, \vec{x}_j) \wedge \bigwedge_{0 \leq j < i \leq n} \neg \varphi(\vec{x}_i, \vec{x}_j) \right).$$

If  $T = \bigcup_{k \in \omega} T_k$ , then these sentences are all in some  $\mathcal{S}_k$ , and then (assuming  $T_k$  is complete)  $T_k$  will be unstable.

An arbitrary theory  $T$  is **companionable** if, for some theory  $T^*$  of its signature,

- $T_{\forall} = T^*_{\forall}$ ,
- $T$  is model-complete.

In this case,  $T^*$  is the **model-companion** of  $T$ . If

$$T_0 \subseteq T_1 \subseteq T_2 \subseteq \dots,$$

and each  $T_k$  has the model-companion  $T_k^*$ , and

$$T_0^* \subseteq T_1^* \subseteq T_2^* \subseteq \dots, \quad (*)$$

then  $\bigcup_{k \in \omega} T_k^*$  is the model-companion of  $\bigcup_{k \in \omega} T_k$ . However (\*) may fail.

**Theorem** (McGrail).  *$m$ -DF<sub>0</sub> (in characteristic 0) has a model-companion,  $m$ -DCF<sub>0</sub>, which admits quantifier-elimination and is  $\omega$ -stable.*

**Theorem** (P.).  *$m$ -DF has a model-companion,  $m$ -DCF. Nevertheless,  $\bigcup_{m \in \omega} m$ -DF is not companionable.*

For the last part (non-companionability), if  $j \in \omega$ , let  $K_j$  be an e.c. (existentially closed) model of  $\omega$ -DF (that is,  $\bigcup_{m \in \omega} m$ -DF), with

$$\mathbb{F}_p(\alpha) \subseteq K_j, \quad \alpha \notin \mathbb{F}_p^{\text{alg}}, \quad \partial_i \alpha = \delta_{ij} = \begin{cases} 1, & \text{if } i = j, \\ 0, & \text{if } i \neq j. \end{cases}$$

Then  $\alpha$  has no  $p$ -th root in  $K_j$ , since

$$\partial_j \alpha = 1, \quad \partial_j(x^p) = p \cdot x^{p-1} \cdot \partial_j x = 0.$$

Therefore  $\alpha$  has no  $p$ -th root in a nonprincipal ultraproduct

$$\prod_{j \in \omega} K_j / \mathfrak{p},$$

even though, in this,  $\partial_i \alpha = 0$  for all  $i$  in  $\omega$ , so  $\alpha$  has a  $p$ -th root in some extension. Thus the ultraproduct is not e.c.. Therefore the class of e.c. models of  $\omega$ -DF is not elementary.

**Theorem (P.).**

$$m\text{-DCF}_0 \subseteq (m+1)\text{-DCF}_0,$$

and therefore  $\omega\text{-DF}_0$  has a model-companion,  $\omega\text{-DCF}_0$ , which is stable, but not superstable.

This is established by means of:

**Theorem (folklore, P.).** *Assuming  $T_0 \subseteq T_1$ , each  $T_k$  having signature  $\mathcal{S}_k$ , consider:*

A. *If*

$$\mathfrak{A} \models T_1, \quad \mathfrak{B} \models T_0, \quad \mathfrak{A} \upharpoonright \mathcal{S}_0 \subseteq \mathfrak{B},$$

*then there is  $\mathfrak{C}$  such that*

$$\mathfrak{C} \models T_1, \quad \mathfrak{A} \subseteq \mathfrak{C}, \quad \mathfrak{B} \subseteq \mathfrak{C} \upharpoonright \mathcal{S}_0.$$

B. *For all  $\mathfrak{A}$ ,*

$$\mathfrak{A} \models_{\text{ec}} T_1 \implies \mathfrak{A} \upharpoonright \mathcal{S}_0 \models_{\text{ec}} T_0.$$

C.  *$T_0$  has the Amalgamation Property: if one model embeds in two others, then those two in turn embed in a fourth model, compatibly with the original embeddings.*

D.  *$T_1$  is  $\forall\exists$  (so that every model embeds in an e.c. model).*

*We have the two implications*

$$A \implies B, \quad B \ \& \ C \ \& \ D \implies A,$$

*but there is no implication among the four conditions that does not follow from these. This is true, even if  $T_1$  is required to be a conservative extension of  $T_0$ , so that  $T_1 \upharpoonright \mathcal{S}_0 = T_0$ .*

*Proof.* (Can be left as exercise.) Suppose A holds. Let

$$\mathfrak{A} \models_{\text{ec}} T_1, \quad \mathfrak{B} \models T_0, \quad \mathfrak{A} \upharpoonright \mathcal{S}_0 \subseteq \mathfrak{B}.$$

We show

$$\mathfrak{A} \upharpoonright \mathcal{S}_0 \preceq_1 \mathfrak{B}$$

(i.e. every existential formula over  $\mathfrak{A} \upharpoonright \mathcal{S}_0$  soluble in  $\mathfrak{B}$  is soluble in  $\mathfrak{A} \upharpoonright \mathcal{S}_0$ ). By hypothesis, there is a model  $\mathfrak{C}$  of  $T_1$  such that

$$\mathfrak{A} \subseteq \mathfrak{C}, \quad \mathfrak{B} \subseteq \mathfrak{C} \upharpoonright \mathcal{S}_0.$$

Then

$$\begin{aligned} \mathfrak{A} &\preceq_1 \mathfrak{C}, \\ \mathfrak{A} \upharpoonright \mathcal{S}_0 &\preceq_1 \mathfrak{C} \upharpoonright \mathcal{S}_0, \\ \mathfrak{A} \upharpoonright \mathcal{S}_0 &\preceq_1 \mathfrak{B}. \end{aligned}$$

Therefore  $\mathfrak{A} \upharpoonright \mathcal{S}_0$  must be an e.c. model of  $T_0$ . Thus  $B$  holds.

Suppose conversely  $B$  &  $C$  &  $D$  holds. Let

$$\mathfrak{A} \models T_1, \quad \mathfrak{B} \models T_0, \quad \mathfrak{A} \upharpoonright \mathcal{S}_0 \subseteq \mathfrak{B}.$$

We establish the consistency of

$$T_1 \cup \text{diag}(\mathfrak{A}) \cup \text{diag}(\mathfrak{B}).$$

It is enough to show the consistency of

$$T_1 \cup \text{diag}(\mathfrak{A}) \cup \{\exists \vec{x} \varphi(\vec{x})\},$$

where  $\varphi$  is an arbitrary quantifier-free formula of  $\mathcal{S}_0(A)$  such that

$$\mathfrak{B} \models \exists \vec{x} \varphi(\vec{x}).$$

By  $D$ , there is  $\mathfrak{C}$  such that

$$\mathfrak{C} \models_{\text{ec}} T_1, \quad \mathfrak{A} \subseteq \mathfrak{C}.$$

By  $B$  then,

$$\mathfrak{C} \upharpoonright \mathcal{S}_0 \models_{\text{ec}} T_0, \quad \mathfrak{A} \upharpoonright \mathcal{S}_0 \subseteq \mathfrak{C} \upharpoonright \mathcal{S}_0.$$

By  $C$ , both  $\mathfrak{B}$  and  $\mathfrak{C} \upharpoonright \mathcal{S}_0$  embed over  $\mathfrak{A} \upharpoonright \mathcal{S}_0$  in a model  $\mathfrak{D}$  of  $T_0$ . In particular,

$$\mathfrak{D} \models \exists \vec{x} \varphi(\vec{x}).$$

Therefore  $\varphi$  is already soluble in  $\mathfrak{C} \upharpoonright \mathcal{S}_0$  itself. Thus

$$\mathfrak{C} \models T_1 \cup \text{diag}(\mathfrak{A}) \cup \{\exists \vec{x} \varphi(\vec{x})\}.$$

Therefore  $A$  holds.

For the rest, 11 (counter-)examples are found. . . □

Now suppose

$$\begin{aligned} (L, \partial_0, \dots, \partial_{m-1}) &\models m\text{-DF}_0, \\ K &\subseteq L, \\ (K, \partial_0 \upharpoonright K, \dots, \partial_{m-1} \upharpoonright K, \partial_m) &\models (m+1)\text{-DF}_0, \\ a &\in L \setminus K. \end{aligned}$$

We shall define a differential field

$$(K\langle a \rangle, \tilde{\partial}_0, \dots, \tilde{\partial}_m),$$

where  $a \in K\langle a \rangle$ , and for each  $i$  in  $m$ ,

$$\tilde{\partial}_i \upharpoonright K\langle a \rangle \cap L = \partial_i \upharpoonright K\langle a \rangle \cap L, \quad (\dagger)$$

and  $\tilde{\partial}_m \upharpoonright K = \partial_m$ .

Considering  $\omega^{m+1}$  as the set of  $(m+1)$ -tuples of natural numbers, we shall have

$$K\langle a \rangle = K(a^\sigma : \sigma \in \omega^{m+1}),$$

where

$$a^\sigma = \tilde{\partial}_0^{\sigma(0)} \dots \tilde{\partial}_m^{\sigma(m)} a. \quad (\ddagger)$$

In particular then, by  $(\dagger)$ , we must have

$$\sigma(m) = 0 \implies a^\sigma = \partial_0^{\sigma(0)} \dots \partial_{m-1}^{\sigma(m-1)} a. \quad (\S)$$

Using this rule, we make the definition

$$K_1 = K(a^\sigma : \sigma(m) = 0).$$

Recursively, we can define

$$K_j = K(a^\sigma : \sigma(m) < j)$$

as desired. If  $L \setminus K\langle a \rangle \neq \emptyset$ , we can repeat, as necessary.

It may not be possible to make  $L$  itself closed under  $\tilde{\partial}_m$ .